

# TxnSails: Achieving Serializable Transaction Scheduling with Self-Adaptive Isolation Level Selection

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# Outline



- Background
- System Overview
- Technique Details
- Evaluations
- Conclusion

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# Background



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- Applications rely databases to store and manage their data.



Serializable, Snapshot isolation, Read committed, which isolation level should I choose ?

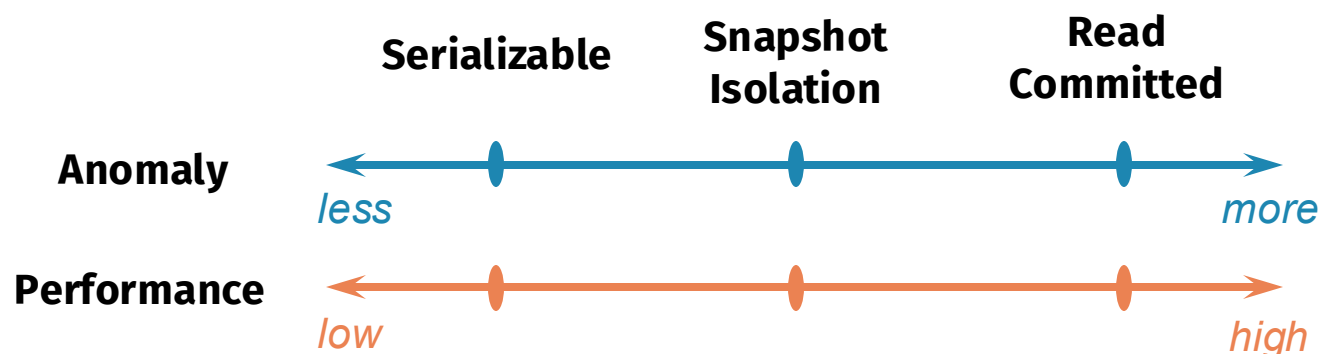
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# Background



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I need both **high performance** and **no data anomaly**,  
what can I do without kernel modification ?



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Configure the database to a low isolation level.

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How to avoid data anomalies when the database  
is configured to low isolation levels ?



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How to avoid data anomalies when the database  
is configured to low isolation levels ?

I. Do concurrency control inside the application <sup>[1]</sup>.

High Performance **but unsafe, easy to write critical bugs**<sup>[1]</sup>

[1] Ad Hoc Transactions in Web Applications: The Good, the Bad, and the Ugly

# Background



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I need both **high performance** and **no data anomaly**,  
what can I do without kernel modification ?

How to avoid data anomalies when the database  
is configured to low isolation levels ?

- I. Do concurrency control inside the application <sup>[1]</sup>.
- II. Promote some specific reads to writes<sup>[2,3]</sup>.

High Performance **but overclaimed, prone to false positives**

[1] Ad Hoc Transactions in Web Applications: The Good, the Bad, and the Ugly

[2] Serializable Use of Read Committed Isolation Level

[3] Robustness against Read Committed for Transaction Templates

# Challenges



## Challenge1:

How to elevate low isolation levels to SER without additional writes ?

# Challenges



## Challenge1:

How to elevate low isolation levels to SER without additional writes ?



Which low isolation level should I choose?

# Challenges



## Challenge1:

How to elevate low isolation levels to SER without additional writes ?

## Challenge2:

How to determine the optimal isolation level for specific workload ?

# Challenges



## Challenge1:

How to elevate low isolation levels to SER without additional writes ?

## Challenge2:

How to determine the optimal isolation level for specific workload ?

## Challenge3:

How to guarantee SER during the transition between isolation levels ?

# Background



## Data Anomaly Structure



## Vulnerable dependency



## Necessary condition of anomaly

Under SI, a **dangerous structure** is defined as the RW-dependency chain,  $T_i \xrightarrow{rw} T_j \xrightarrow{rw} T_k$ . In a data anomaly, at least one dangerous structure exists, and  $T_k$  commits before  $T_j$  commits.

Under RC, a single RW-dependency,  $T_i \xrightarrow{rw} T_j$ , is referred to as a **dangerous structure**. In a data anomaly, at least one dangerous structure exists, and  $T_j$  commits before  $T_i$  commits.

1. The RW dependency  $T_j \xrightarrow{rw} T_k$  in chain  $T_i \xrightarrow{rw} T_j \xrightarrow{rw} T_k$  under SI.
2. The single RW dependency  $T_i \xrightarrow{rw} T_j$  under RC.

1. A vulnerable dependency exists.
2. The **dependency order** between the two transactions in the vulnerable dependency is inconsistent with their **commit order**.

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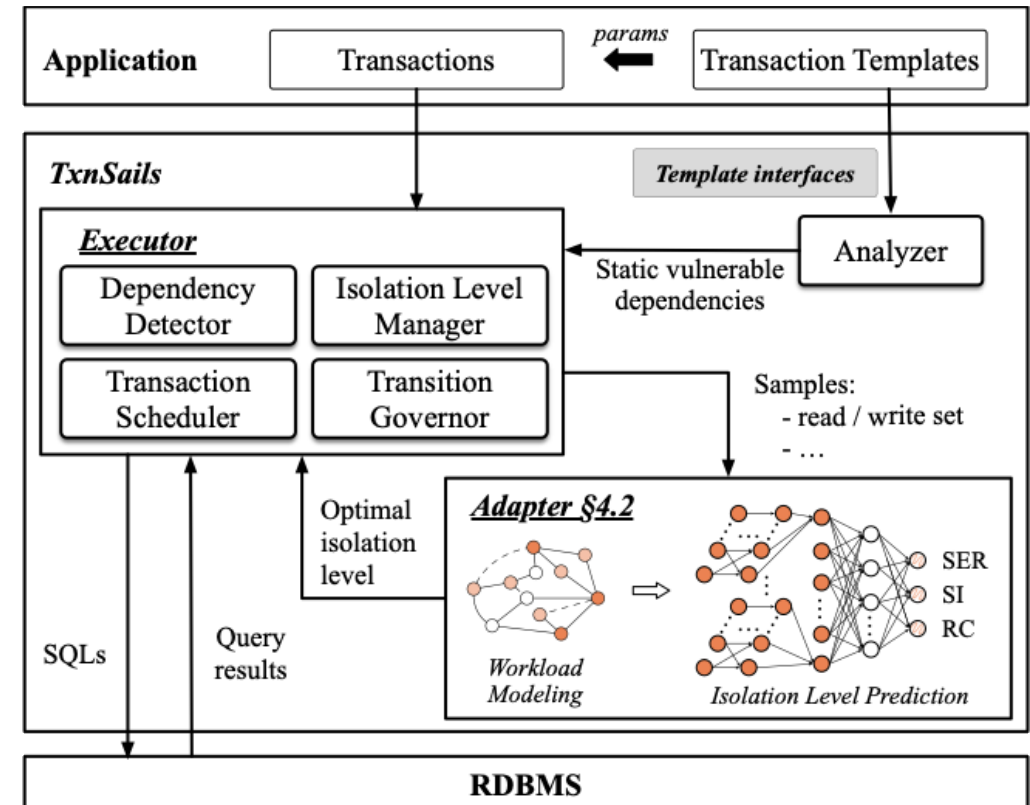


# System Overview



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TxnSails works in the middle tier between the application tier and database tier, it comprises three main components: ***Analyzer***, ***Executor***, and ***Adapter***.



# System Overview

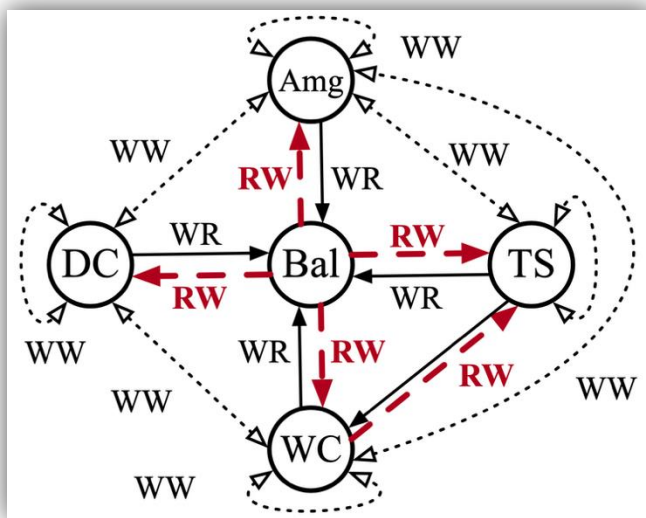


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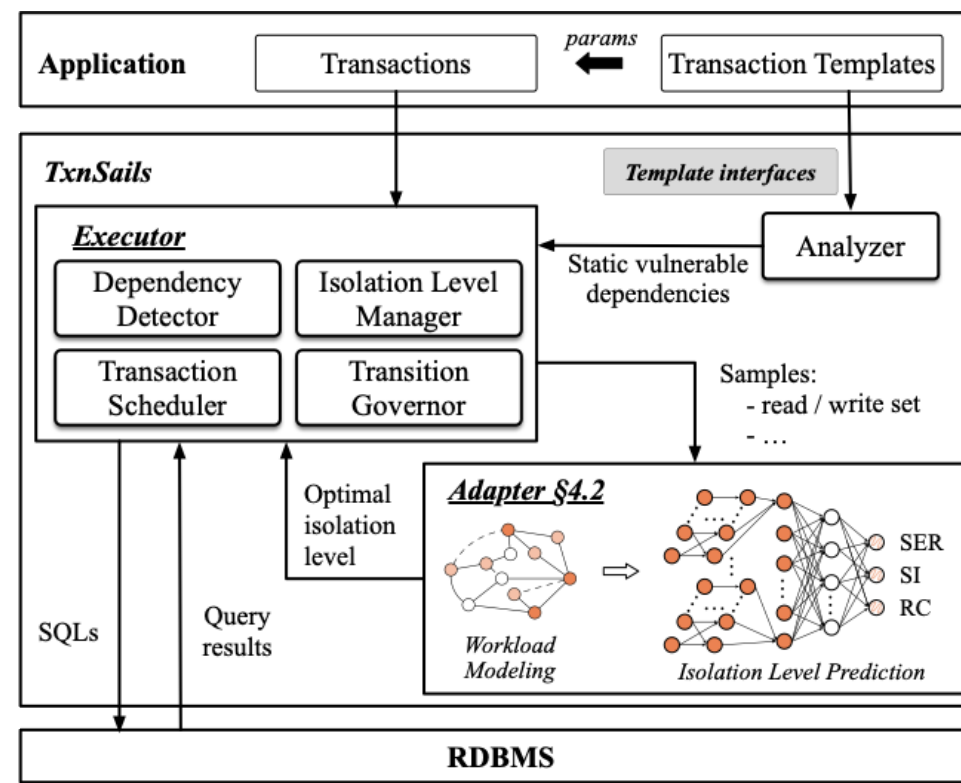
TxnSails works in the middle tier between the application and database, it comprises three main components: **Analyzer**, **Executor**, and **Adapter**.

## ✂ Analyzer

It builds the static dependency graph (SDG) for the transaction templates and identifies all the static vulnerable dependencies for each low isolation level.



Smallbank SDG



# System Overview



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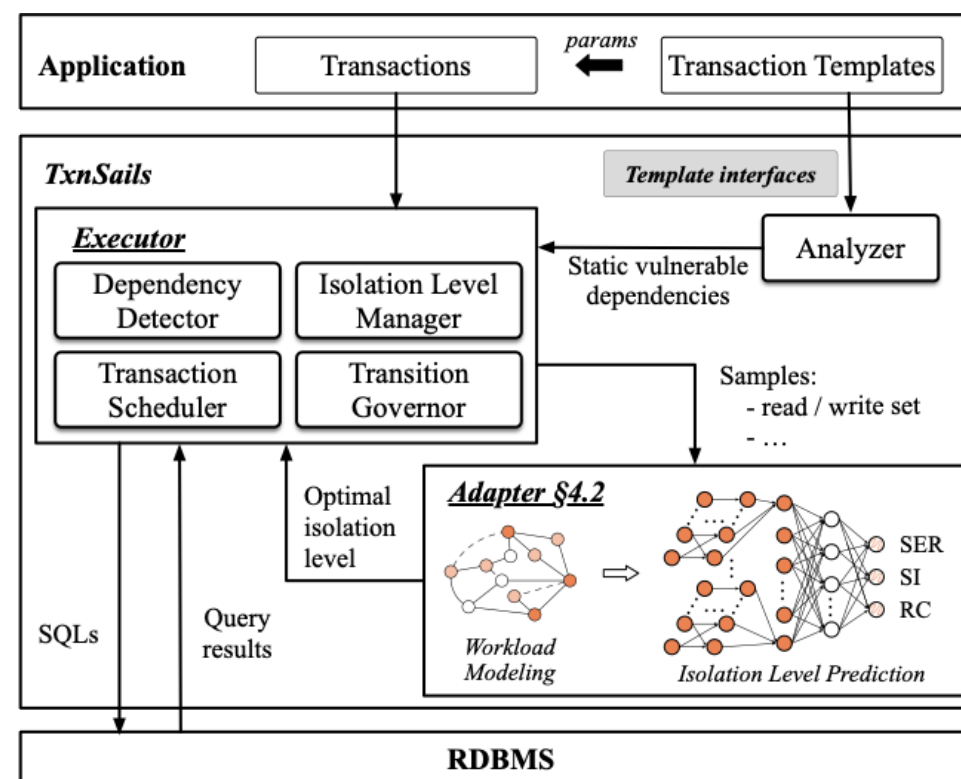
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## ✂ Analyzer

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## ✂ Executor

It ensures SER when transactions operate at a single low isolation level or during the transition. Keep the commit order consistent with the dependency.



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## ✂ Analyzer

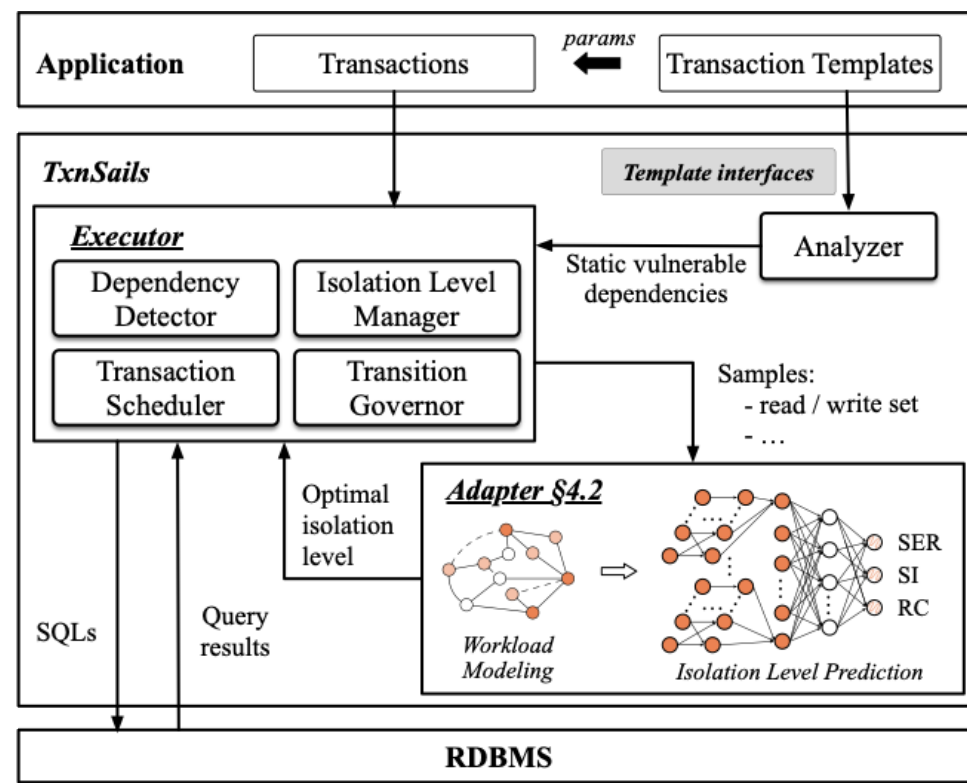
It builds the static dependency graph (SDG) for the transaction templates and identifies all the static vulnerable dependencies for each low isolation level.

## ✂ Executor

It ensures SER when transactions operate at a single low isolation level or during the transition. Keep the commit order consistent with the dependency.

## ✂ Adaptor

It collects transactions and then employs a graph-based model to predict the optimal isolation level.



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# Technique Details

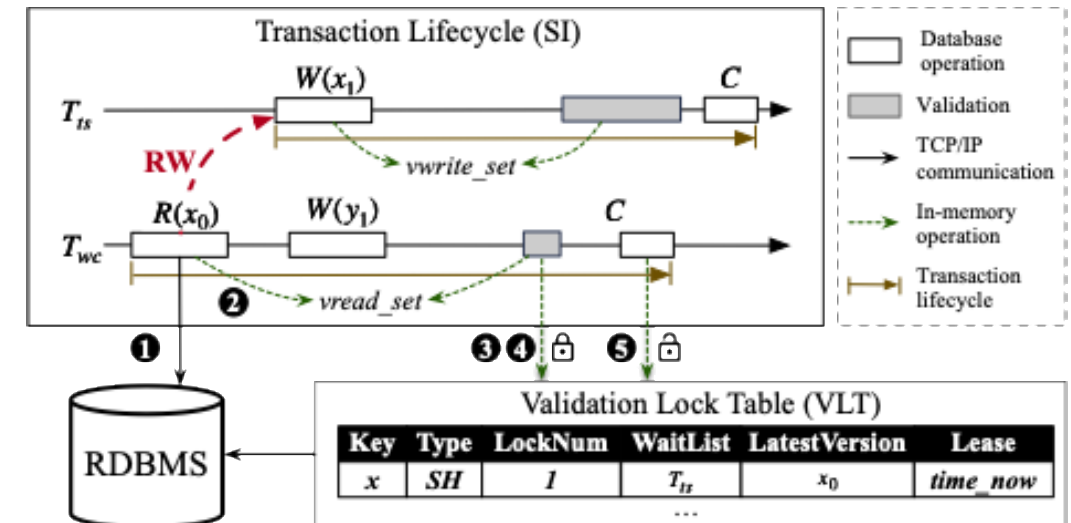


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## ✚ Unified middle-tier concurrency control mechanism

TxnSails dynamically detects runtime vulnerable dependencies and schedules their commit order. Specifically, it adds a validation phase in transaction lifecycle before these transactions can commit.

- I. In the execution phase, TxnSails stores the read/write data items in the thread-local buffer that may induce the vulnerable dependencies;
- II. In the validation phase, TxnSails acquires validation locks for data items stored in the buffer. Then, it detects dependencies among them and aims to schedule the commit order consistent with the identified dependency order;
- III. In the commit phase, TxnSails applies modifications to the database and subsequently releases the validation locks.



# Technique Details



## Self-adaptive isolation level selection

TxnSails adopts transaction dependency graphs to capture workload features and adopts a graph classification model to perform self-adaptive isolation level selection.

### **Graph construction**

TxnSails proposes a graph-structured workload model, where each transaction is mapped to a vertex  $v_i$ , and its feature vector  $v_i$  is generated by extracting the number of data items in its read and write set. For each edge  $(v_i, v_j)$ , TxnSails extracts the data dependency type and the involved relations to generate its feature.

### **Graph embedding and isolation level prediction**

The graph model comprises two parts. First, we use a *Graph Embedding Network* to aggregate both vertex and edge features, encoding the local structure and attribute information of the graph. Second, we use a *Graph Classification Network* that learns the mapping from the embedded matrix to the optimal isolation level and perform the graph classification to predict the optimal isolation level.



# Technique Details



## Cross-isolation level validation mechanism

If the predicted optimal isolation level changes, TxnSails will adapt from the previous isolation level  $I_{old}$  to the optimal isolation level  $I_{new}$ . TxnSails employs a cross-isolation validation (CIV) mechanism that ensures serializability and allows for non-blocking transaction execution.

### **Cross-isolation vulnerable dependency**

The cross-isolation vulnerable dependency is defined as  $T_j \xrightarrow{rw} T_k$  in chain  $T_i \xrightarrow{rw} T_j \xrightarrow{rw} T_k$ , where  $T_j$  commits after the transition starts.



### **Corollary**

For any cross-isolation vulnerable dependency  $T_j \xrightarrow{rw} T_k$ , if  $T_j$  commits before  $T_k$ , the scheduling achieves SER.



# Technique Details



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### **Corollary**

For any cross-isolation vulnerable dependency  $T_j \xrightarrow{rw} T_k$ , if  $T_j$  commits before  $T_k$ , the scheduling achieves SER.

**Do middle-tier concurrency control to ensure the consistent dependency order and commit order !**

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# Evaluations

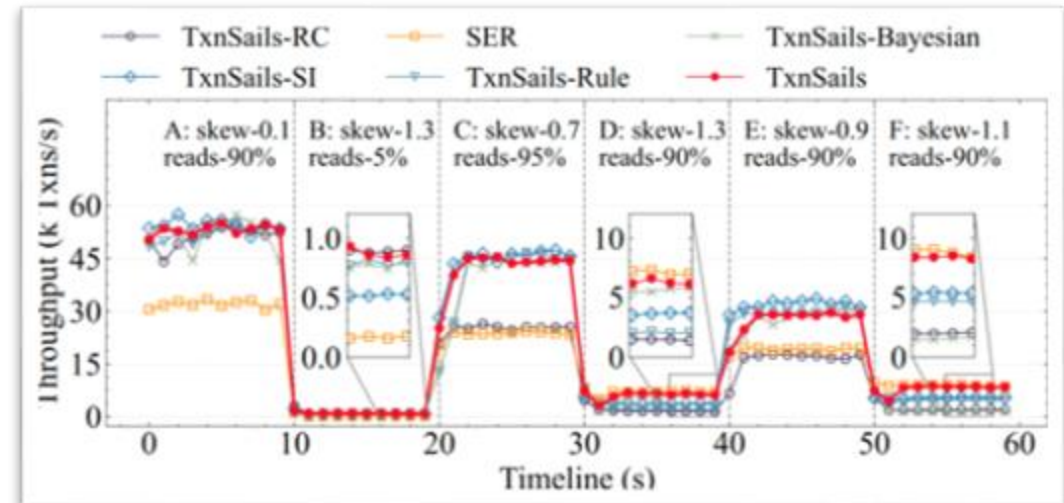


## ★ Self-adaptive isolation level selection

We first evaluate selfadaptive isolation level selection by varying the workload every 10s across six distinct scenarios. We sample the workload at 1-second intervals. The results demonstrate that different isolation levels perform variably under different workloads:

- ✓ SI performs well in low-skew scenarios (A,C,E).
- ✓ SER is more suitable in high-skew scenarios with little writes (D,F).
- ✓ RC excels in high skew scenarios with more writes (B).

***TxnSails can choose the optimal isolation level across all tested scenarios.***



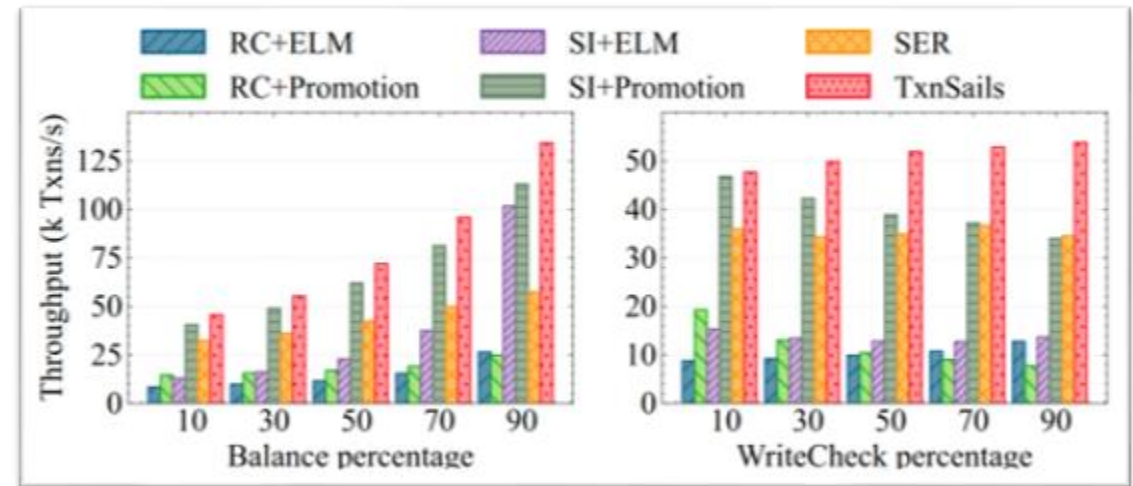
# Evaluations



## ★ Impact of templates percentages

In complex workloads like SmallBank and TPC-C, only certain queries lead to data anomalies. This part compares different solutions by varying the percentage of critical transaction templates.

As the ratio of Balance transaction increases, TxnSails achieves up to 6.2× performance gain. As the ratio of WriteCheck transactions increases, TxnSails's advantage becomes more pronounced as no extra WW conflicts are introduced, outperforming up to 2.3× the performance of SER.



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# Conclusion



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I will choose TxnSails to manage database isolation level.

- TxnSails introduces a unified middle-tier validation method to enforce the commit order consistent with the vulnerable dependency order, ensuring serializability in single-isolation and cross-isolation scenarios.
- TxnSails adopts a graph learned model to extract the runtime workload characteristics and adaptively predict the optimal isolation levels, achieving further performance improvement.

# Thanks!

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